Pathfinding in 3D Space

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Outline

- Introduction
- ► I. State of the art
- ► II. Algorithms
- III. Implementation in 3D space
- ► IV. Results
- Conclusion

Introduction

Objective: Find the shortest paths efficiently in 3D space

Applications: video games, drone navigation

► Homeworld (1999) :

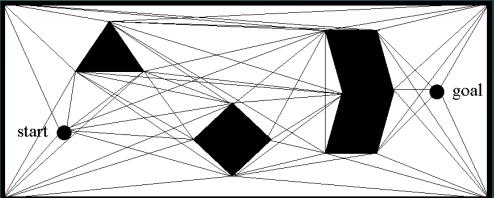
First famous real-time strategy game with movement in 3D space



Shortest paths in a graph

- Dijkstra (single source)
 - ► O((|V|+|E|)log(|V|))
- Bellman-Ford (single source, weighted directed graph)
 - ► O(|V||E|)
- Floyd-Warshall (for all pairs of vertices, weighted graph , no negative cycle)
 O(|V|³)
- A* (single source, single destination)
 - \blacktriangleright O(n), n = length of the solution path => O(|E|)

- 2D exact
 - Visibility graph
 - Anya (2D grid)
- 2D approximate
 - ► Waypoints
 - Navigation mesh + tunnel
 - ► Family of Theta*



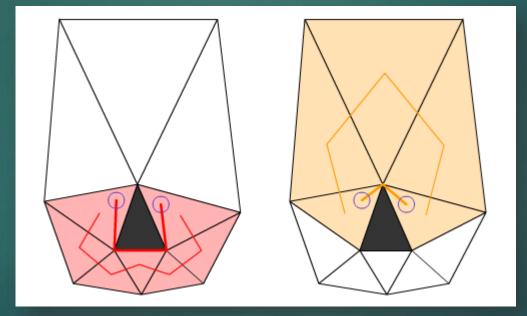






Non-optimality

Navigation mesh + tunnel



path found VS true shortest path

3D surface - exact

- Windows (Fast exact and approximate geodesics on meshes 2005 Surazhsky)
- ► 3D surface approximate
 - Heat (Geodesics in heat 2013 Crane)
 - Fast-marching (1996 Sethian)

- World representation
 - ► Tetrahedralization
 - Convex decomposition
 - ► Grid
 - ► Octree

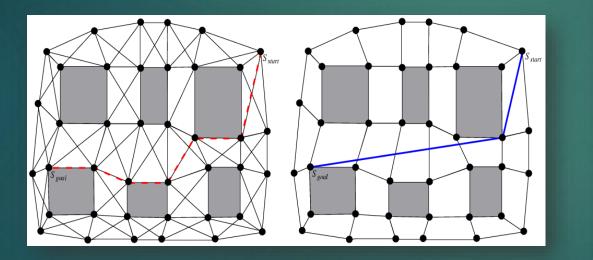
A* (1968 Hart)

h admissible if no over-estimation and h(y) <= h(x) + d(x, y)

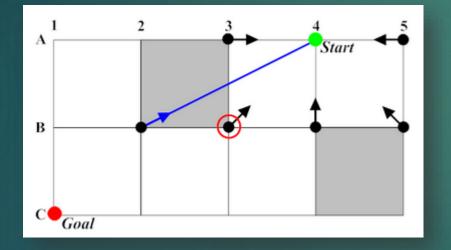


1 Main() $open := closed := \emptyset;$ 2 $g(s_{start}) := 0;$ 3 $parent(s_{start}) := s_{start};$ 4 $open.Insert(s_{start}, g(s_{start}) + h(s_{start}));$ 5 while open $\neq \emptyset$ do 6 s := open.Pop();7 if $s = s_{\text{goal}}$ then 8 return "path found"; Q $closed := closed \cup \{s\};$ 10 foreach $s' \in nghbr_{vis}(s)$ do 11 if $s' \notin closed$ then 12 if $s' \notin open then$ 13 $g(s') := \infty;$ 14 parent(s') := NULL;15 UpdateVertex(s, s'); 16 return "no path found"; 17 18 end UpdateVertex(s, s') 19 $q_{old} := q(s');$ 20 ComputeCost(s, s'); 21 if $g(s') < g_{old}$ then 22 if $s' \in \text{open then}$ 23 open.Remove(s'); 24 open.Insert(s', g(s') + h(s'));25 26 end ComputeCost(s, s') 27 /* Path 1 */ 28 if g(s) + c(s, s') < g(s') then 29 parent(s') := s;30 g(s') := g(s) + c(s, s');31 32 end

► Theta* (2007 Nash)



Lazy Theta* (2010 Nash)

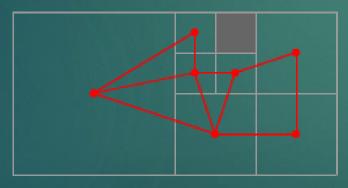


| | F • 0 |
|------------|---|
| | Aain() |
| 2 | $open := closed := \emptyset;$ |
| 3 | $g(s_{start}) := 0;$ |
| 4 | $parent(s_{start}) := s_{start};$ |
| 5 | $open.Insert(s_{start}, g(s_{start}) + h(s_{start}));$ |
| 6 | while open $\neq \emptyset$ do |
| 7 | s := open.Pop(); |
| 8 | SetVertex (s) ; |
| 9 | if $s = s_{\text{goal}}$ then |
| 10 | return "path found"; |
| 11 | $closed := closed \cup \{s\};$ |
| 12 | foreach $s' \in nghbr_{vis}(s)$ do |
| 13 | if $s' \notin closed$ then |
| 14 | if $s' \not\in$ open then |
| 15 | $g(s') := \infty;$ |
| 16 | parent(s') := NULL; |
| 17 | UpdateVertex (s, s') ; |
| | |
| 18 | return "no path found"; |
| 19 e | nd |
| | JpdateVertex(s, s') |
| 20 U 21 | $ g_{old} := g(s');$ |
| 21 | $g_{old} := g(s),$ ComputeCost $(s, s');$ |
| 22 | if $g(s') < g_{old}$ then |
| 23 24 | $ \mathbf{if} g(s) \rangle \leq g_{old} \mathbf{then}$ |
| 24 | open.Remove(s'); |
| | |
| 26 | open.Insert $(s', g(s') + h(s'));$ |
| 27 e | nd |
| 28 (| ComputeCost(s, s') |
| 29 | /* Path 2 */ |
| 30 | if $g(parent(s)) + c(parent(s), s') < g(s')$ then |
| 31 | parent(s') := parent(s); |
| 32 | g(s') := g(parent(s)) + c(parent(s), s'); |
| | |
| 33 e | nd |
| 34 S | etVertex(s) |
| 35 | if NOT lineofsight(parent(s), s) then |
| 36 | /* Path 1*/ |
| 37 | parent(s) := |
| | $argmin_{s' \in nghbr_{vis}(s) \cap closed}(s(s') + c(s', s));$ |
| 38 | g(s) := |
| | $\min_{s' \in nghbr_{vis}(s) \cap closed}(g(s') + c(s', s));$ |
| 20 | |
| 39 e | nd |

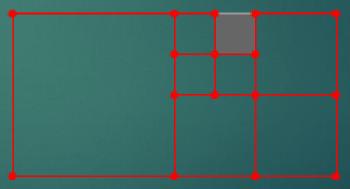
III. Implementation

Octree construction

- ► Triangle-cube intersection
- Progressive octree
- Graph construction



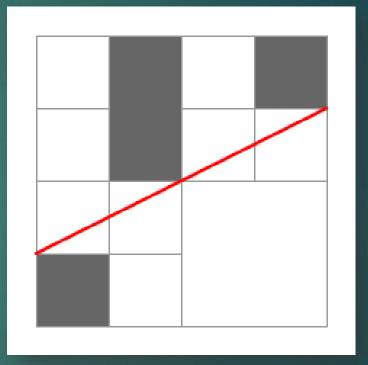
Dual graph (not standard)



Edge-corner

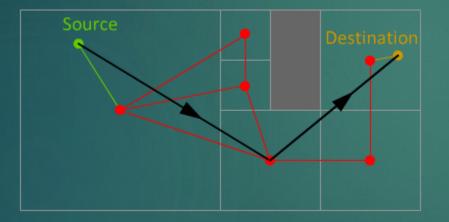
III. Implementation

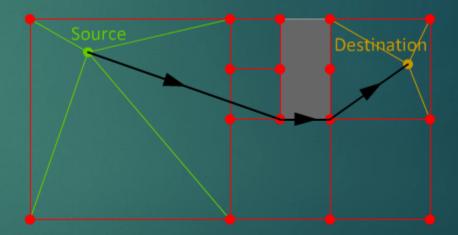
- Line of sight
 - Fast
 - Robust



III. Implementation

Injection of source and destination





III. Implementation - Optimisation

Avoid exhaustive search

Precompute the connectivity of the graph nodes

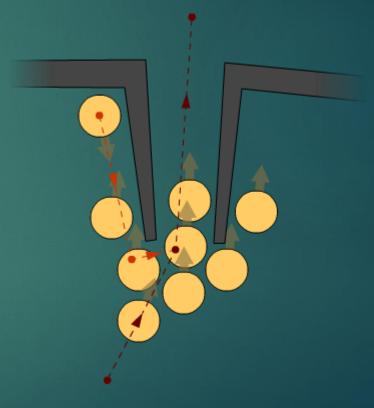
III. Implementation - Optimisation

Multisource

Reuse information

III. Implementation - Extension

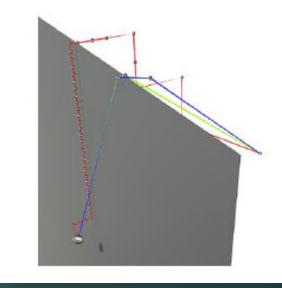
- Application in video games
 - ► Waypoints
 - ► Repulsive force
 - ► Replanning





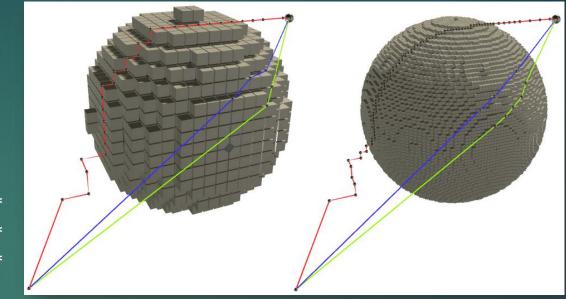
| Data Structure | Algorithm | distance | time cost | distance | time cost |
|--------------------|-------------------------|----------|--------------------|----------|--------------------|
| Octree | A* | 125.18% | $1.6 \mathrm{ms}$ | 124.56% | $4.5 \mathrm{ms}$ |
| | Theta* | 101.58% | $11.8 \mathrm{ms}$ | 109.24% | $53.0 \mathrm{ms}$ |
| | Lazy Theta [*] | 101.73% | $6.6\mathrm{ms}$ | 109.48% | $22.1 \mathrm{ms}$ |
| Progressive Octree | A* | 125.19% | 4.6ms | 126.27% | $5.8 \mathrm{ms}$ |
| | Theta* | 101.58% | $21.5 \mathrm{ms}$ | 101.63% | $42.7 \mathrm{ms}$ |
| | Lazy Theta [*] | 101.39% | $10.6 \mathrm{ms}$ | 102.27% | $18.1 \mathrm{ms}$ |

Table 1: Comparison - A single wall - Left: source sparse/ Right: destination sparse



Red: A* Green: Theta* Blue: Lazy Th<u>eta</u>*

IV. Results



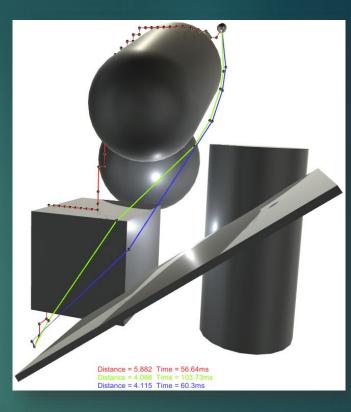
Red: A* Green: Theta* Blue: Lazy Theta*

| Data Structure | Algorithm | distance | time cost | distance | time cost |
|--------------------|-------------------------|----------|-------------------|----------|---------------------|
| Octree | A* | 121.38% | $1.5 \mathrm{ms}$ | 121.03% | $28.0 \mathrm{ms}$ |
| | Theta* | 104.59% | $6.7\mathrm{ms}$ | 102.71% | $236.1 \mathrm{ms}$ |
| | Lazy Theta [*] | 104.65% | $4.2 \mathrm{ms}$ | 102.34% | $114.8 \mathrm{ms}$ |
| Progressive Octree | A* | 127.43% | 3.3ms | 126.88% | $55.6 \mathrm{ms}$ |
| | Theta* | 103.49% | $6.3 \mathrm{ms}$ | 101.23% | $229.4 \mathrm{ms}$ |
| | Lazy Theta [*] | 103.68% | $3.2\mathrm{ms}$ | 101.16% | $108.8 \mathrm{ms}$ |

Table 2: Comparison - A single sphere - Left: level = 7/ Right: level = 9

IV. Results

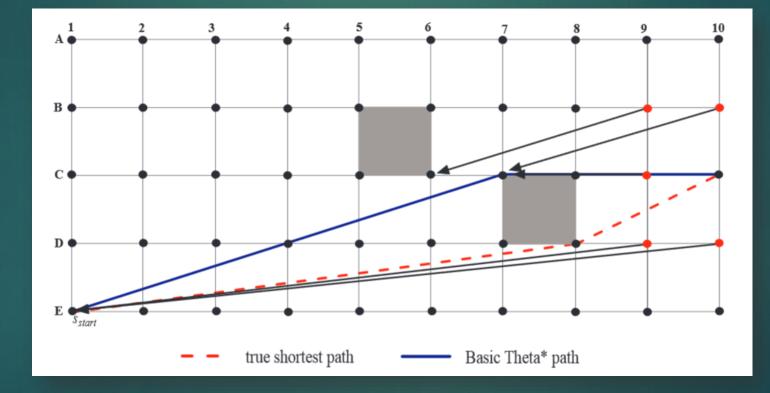
Red: A* Green: Theta* Blue: Lazy Theta*



| Data Structure | Algorithm | distance | time cost | distance | time cost |
|--------------------|-------------------------|----------|---------------------|----------|--------------------|
| Octree | A* | 3.2472 | 9.5ms | 3.3302 | $5.3 \mathrm{ms}$ |
| | Theta* | 2.4108 | $23.9 \mathrm{ms}$ | 2.4600 | $45.5 \mathrm{ms}$ |
| | Lazy Theta [*] | 2.4135 | $9.5\mathrm{ms}$ | 2.4592 | $16.3 \mathrm{ms}$ |
| Progressive Octree | A* | 3.3949 | 14.1ms | 3.3222 | 7.15ms |
| | Theta* | 2.4009 | $17.59 \mathrm{ms}$ | 2.4158 | $43.1 \mathrm{ms}$ |
| | Lazy Theta [*] | 2.4057 | $7.78 \mathrm{ms}$ | 2.4205 | 14.6ms |

Table 3: Comparison - A complex scene - Left: edge-corner graph/ Right: dual graph

Non-optimality of Theta*



Demo!



- ▶ Demo!
 - ► Demo!
 - Demo !
 - Demo !



Conclusion

- Exploration in a new domain
- Our proposition : Lazy Theta * + Progressive Octree + Edge-corner graph
- Possible Improvements
 - Distribution of computation at each frame
 - Other possibilities of h
 - Post-processing

